

Introduction to Formal Methods

Chapter 11. Abstraction Methods

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11. Abstraction Methods

- Abstraction Methods
 - A family of techniques used to simplify automata
 - Simplification aiming at verifying a system (faster) using a model checking approach
 - Examples:
 - (Pb1) " Does $A \models \phi$? " \leftarrow a complex problem
 - (Pb2) " Does $A' \models \phi'$? " \leftarrow a much simpler problem
 - " tricks of the trade "
- Organization of Chapter 11
 - When Is Model Abstraction Required?
 - Abstraction by State Merging
 - What Can Be Proved in the Abstract Automaton?
 - Abstraction on the Variables
 - Abstraction by Restriction
 - Observer Automata

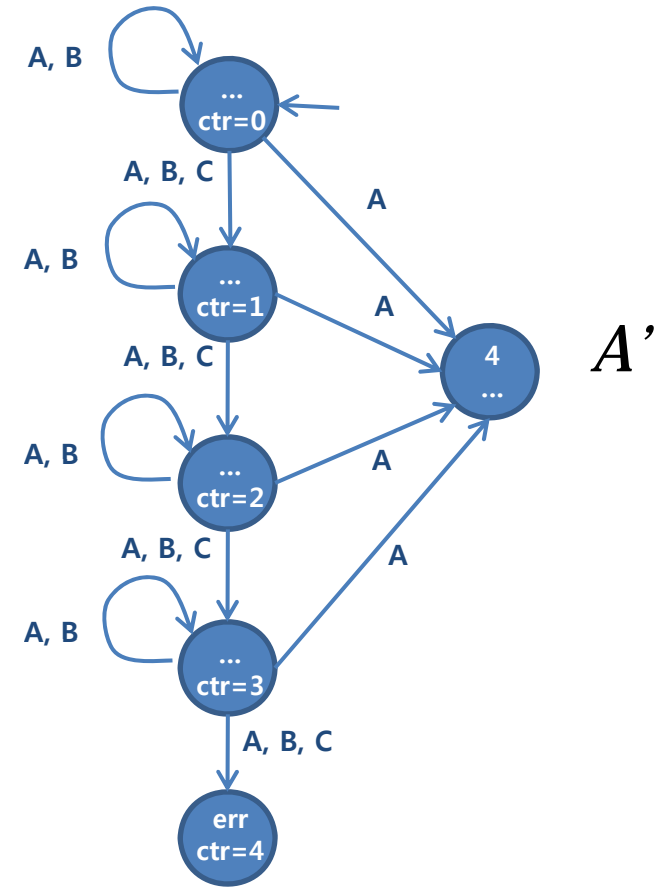
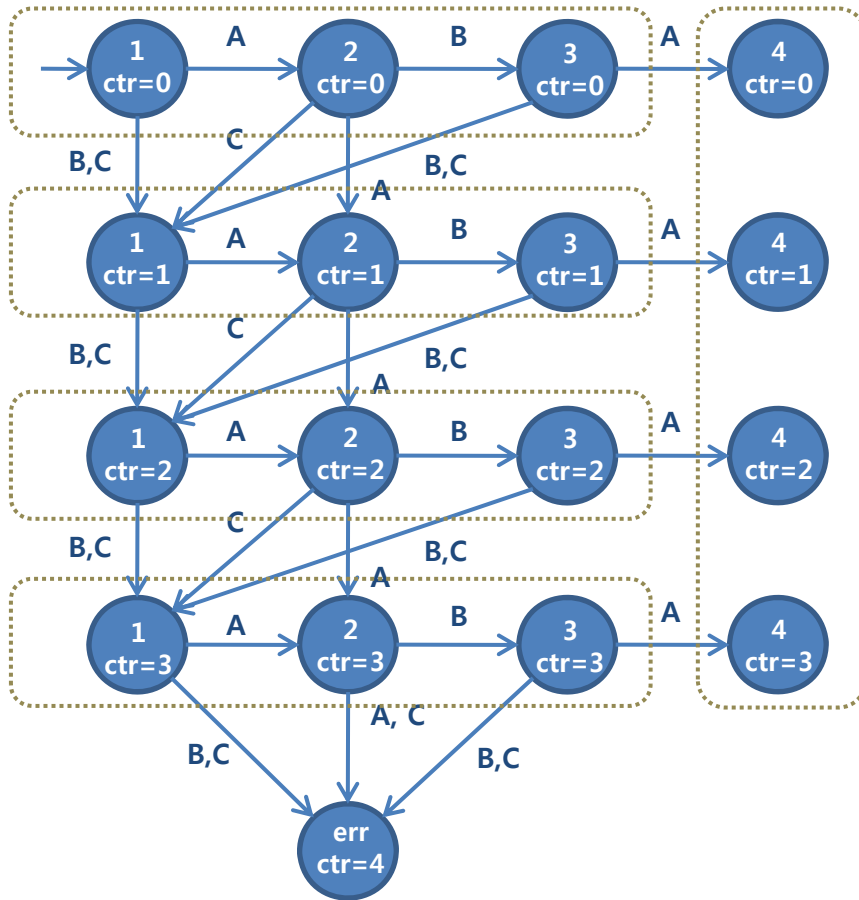
11.1 When Is Model Abstraction Required?

- Two main types of situations for model abstraction
 1. Size of the automaton
 - Too large :
 - Too many variables
 - Too many automata in parallel
 - Too many clocks in the timed automata
 2. Type of the automaton
 - Other types of automata
 - Using integer variables, communication channels, clocks, priorities, etc.
- Three classical abstraction methods
 1. Abstraction by State Merging
 2. Abstraction on the Variables
 3. Abstraction by Restriction

11.2 Abstraction by State Merging

- Folding
 - Viewing some states of an automaton as identical
 - The most important question : Correctness!
 - For example,
 - The digicode door lock with error counters (in Chapter 1)
 - Focusing on the error counter.
 - Correctness problem:
 - All states in A' can be reached through the letter A , but not in A

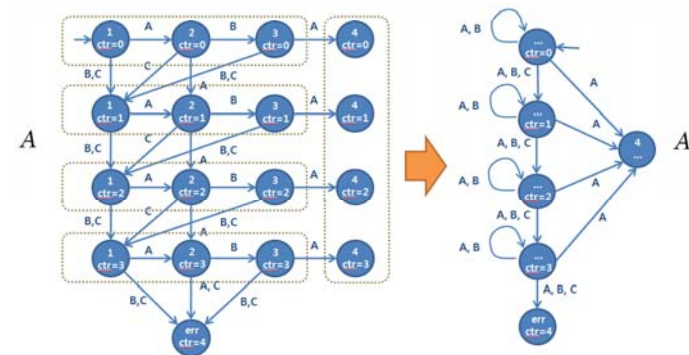
A



A'

11.3 What Can be Proved in the Abstract Automaton?

- We can use state merging to verify safety properties
- Observation (Merging states from A to A')
 1. A' has more behaviors than A .
 2. Now the more behaviors an automaton has, the fewer safety properties it fulfills.
 3. Thus, if A' satisfies a safety property ϕ then a fortiori A satisfies ϕ .
 4. However, if A' does not satisfy ϕ , no conclusions can be drawn about A .
- More behaviors
 - A' has more behaviors than A
 - All executions of A remain present (in folded form) in A'
 - Some new behaviors may be introduced in A'
 - For example, many infinite loops are possible in A'

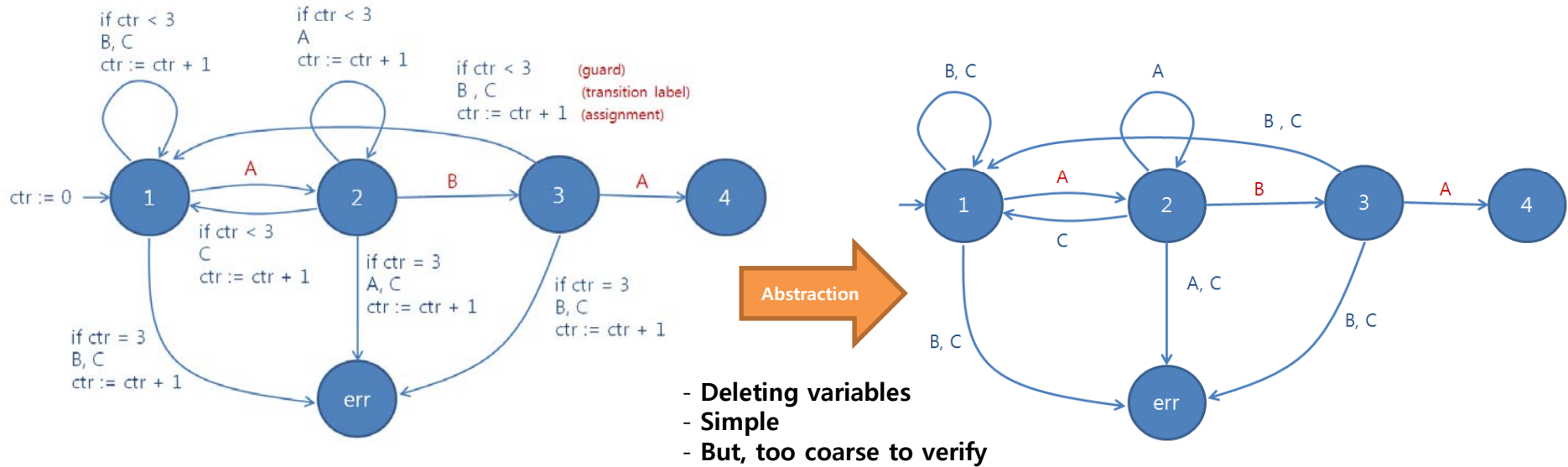


- Preserving safety properties
 - Necessary to ensure that the property ϕ is indeed a safety property.
- One-way preservation
 - If A' does not satisfies ϕ , then A' satisfies $\neg\phi$.
 - But, in general the negation of a safety property is not a safety property.
 - Abstraction methods are often one-way:
 - If the answer is positive, then is positive too.
 - If the answer is negative, then we learned nothing about A .
- Some necessary precautions
 - Skipped.
 - about the propositions' merging and marking in model checking algorithms
- Modularity
 - State merging is preserved by product.
 - $A' \parallel B$ can be obtained from $A \parallel B$ by a merging operation
- State merging in practice
 - Question : " How will we guess and then specify the sets of states to be merged ? "
 - Answer : " The user is the one who defines and applies his own abstraction. "
 - " No tool assistance is offered. "
 - Abstraction on variables are often easy to define and implement.

11.4 Abstraction on the Variables

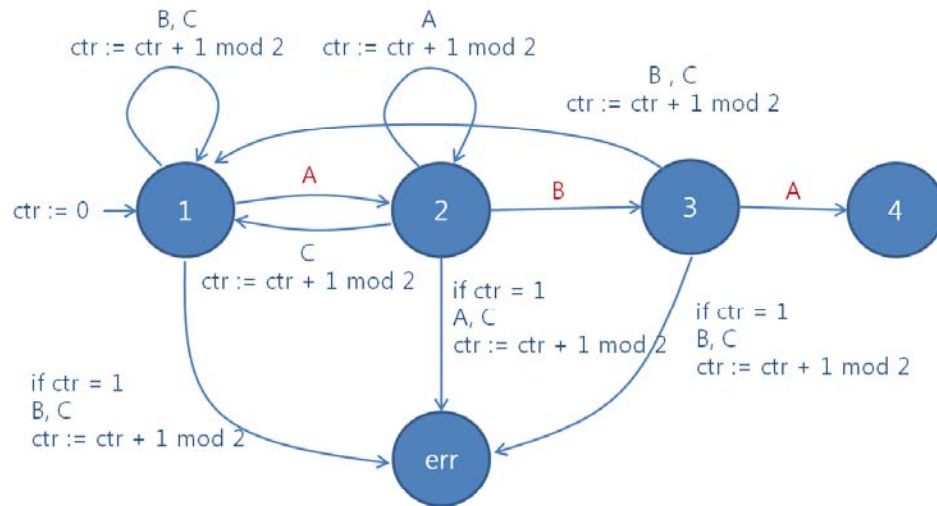
- Abstraction on the variables
 - Concerns the "data" part of automata with variables
 - Directly applies to the description of the automata with variables
- Example

var ctr: int;



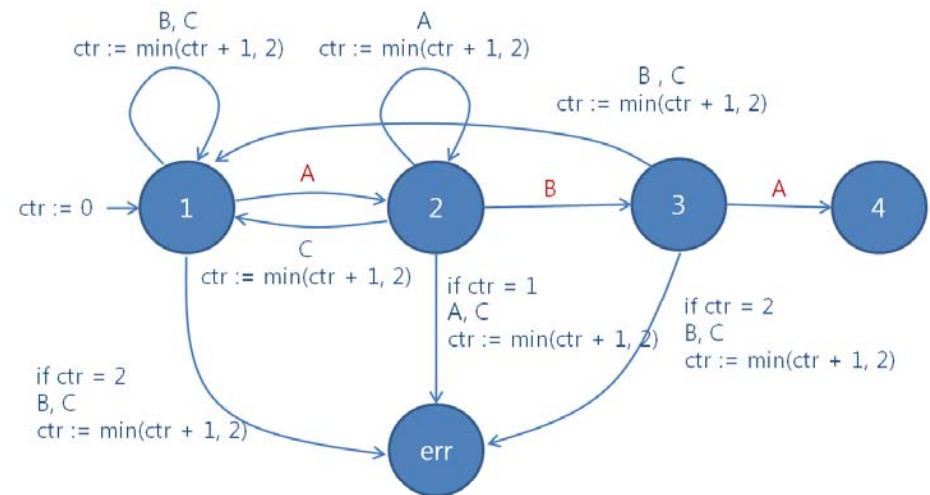
- Abstraction differs from deletion
 - Abstract Interpretation
 - Mathematical theory aiming at defining, analyzing, justifying methods based on abstraction
- Bounded variables
 - Narrow down the domain of variables
 - For example,
 - Integer \rightarrow 0 ~ 10 value
 - The digicode with a modulo 2 counter

var ctr : 0..1



The digicode with a modulo 2 counter

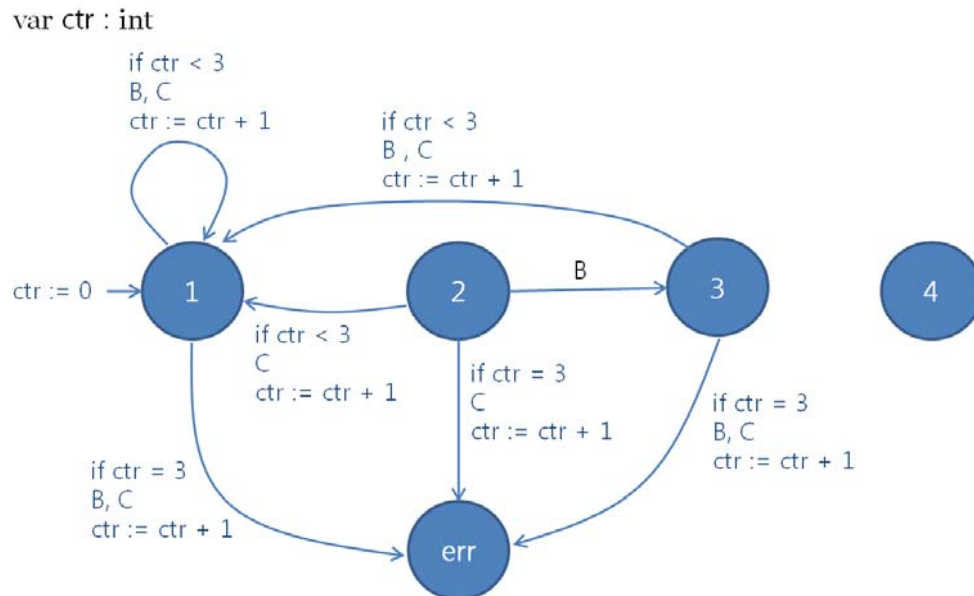
var ctr : 0..2



The digicode with a counter bounded by 2

13.5 Abstraction by Restriction

- Restriction
 - A particular form of simplification
 - Operates by forbidding some behaviors of the system or by making some impossible
 - Removing states or transitions
 - Strengthening the guard, etc.
 - For example
 - Remove all the transitions labeled *A*



The digicode with no *A* transition



The unfolding of the digicode with no *A* transition

- What the restrictions preserve

- If A' is obtained from A by restriction, then literally all the behaviors of A' are behaviors of A .
- Thus if A' does not satisfy a safety property, then a fortiori neither does A .
- Conditional reachability property "EF err" = negation of safety property

- For example,

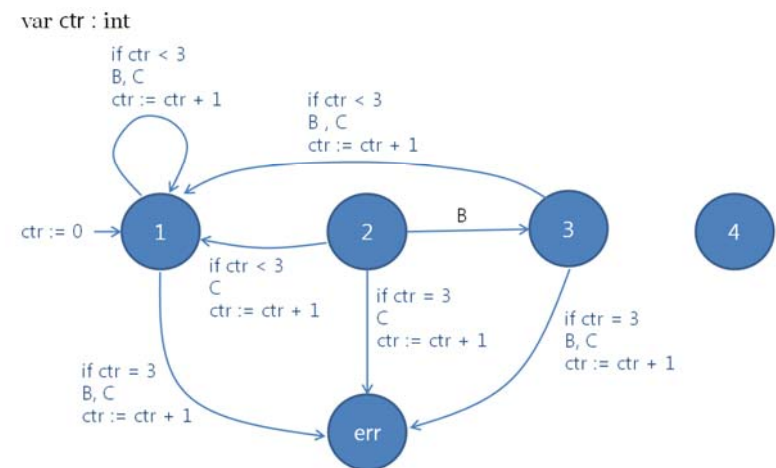
- A' satisfies EF err
- So we conclude that A also satisfies this property

- Inverse preservation

- A safety property does not hold. (To find errors)
- But, not to prove the correctness of A

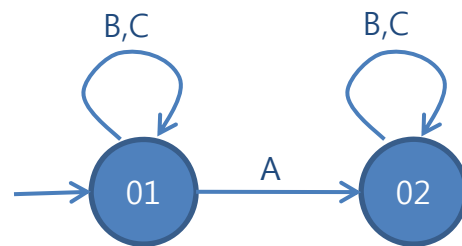
- Advantage of restriction

- Simplicity in conceptual and implementational
- It is a modular operation
- It naturally applies to an automaton with variables

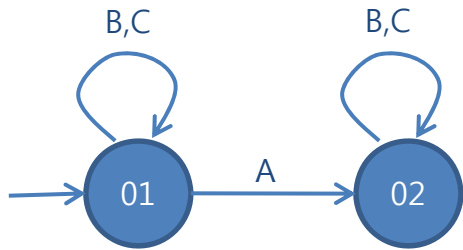


11.6 Observer Automata

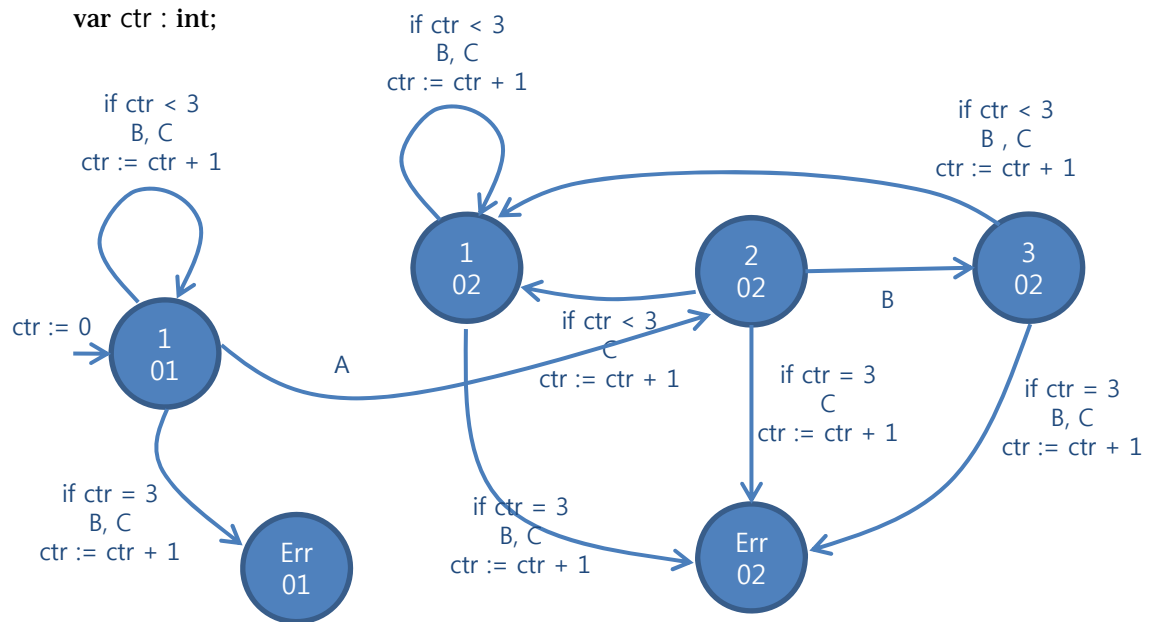
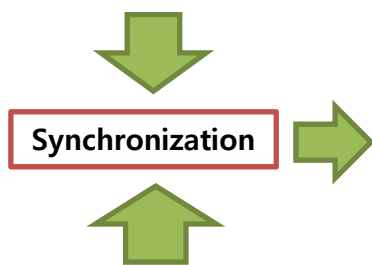
- Observer automata
 - Aiming at simplifying a system by restricting its legitimate behaviors to those accepted by an automata outside the system, called observer automata.
 - Reduce the size of automata by restricting its behavior rather than its structure (states and transitions in restriction methods)
 - PLTL model checking algorithm (in Chapter 3) use the concept.
- An example
 - Supposed that a single A may occur to prove the property.



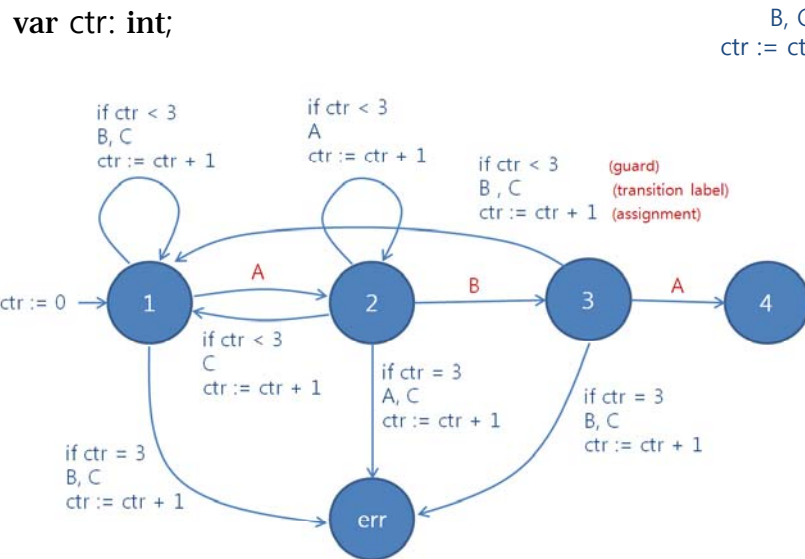
An observer automaton O



An observer automaton O



The synchronized digicode with its observer



An automaton A for the digicode